

A Brief Overview of 802.11 Wireless Networking

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CS6250: Computer Networking

November 10, 2011

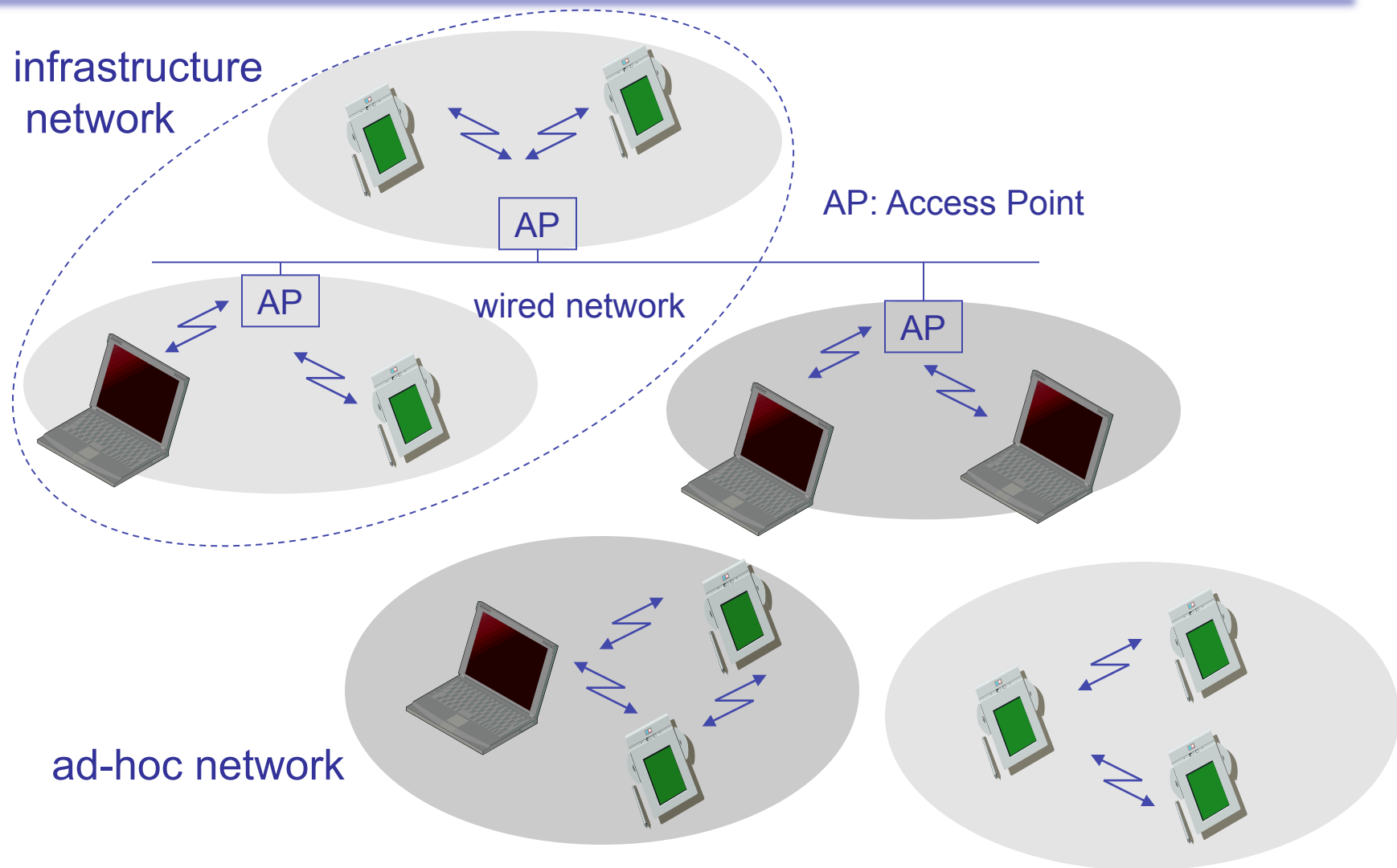
(Some slides curtesy Lili Qiu & Nitin Vaidya)



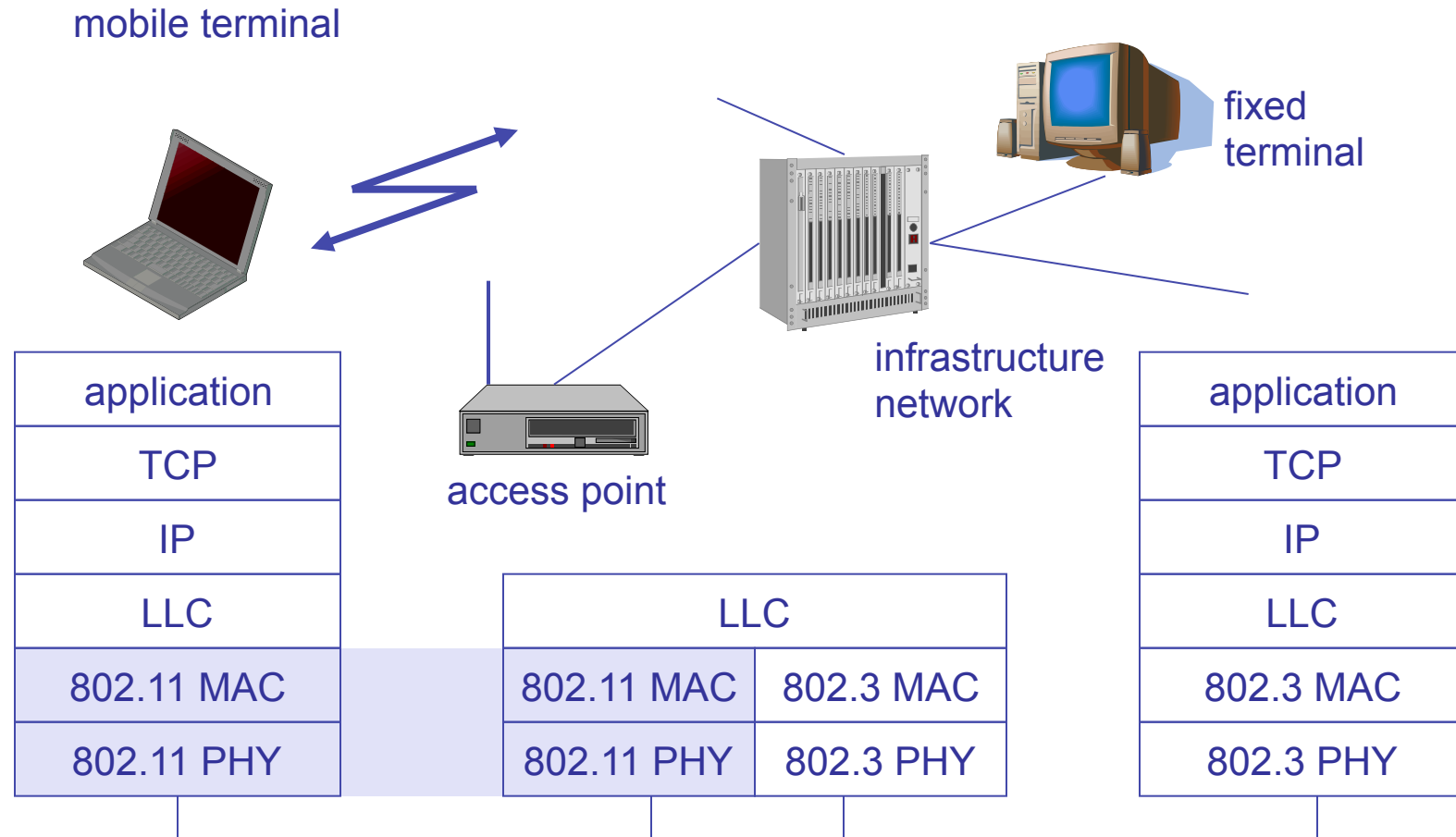
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Infrastructure vs. *Ad hoc*



IEEE 802.11 Infrastructure



802.11 - Layers and functions



- MAC
 - ♦ access mechanisms, fragmentation, error control, encryption
- MAC Management
 - ♦ synchronization, roaming, MIB, power management

- PLCP Physical Layer Convergence Protocol
 - ♦ clear channel assessment signal (carrier sense)
- PMD Physical Medium Dependent
 - ♦ modulation, coding
- PHY Management
 - ♦ channel selection, MIB
- Station Management
 - ♦ coordination of all management functions

DLC	LLC	Station Management
	MAC	
PHY	PLCP	
	PMD	



802.11 Physical Layers

- 802.11b - 2.4 GHz ISM band
 - ◆ FHSS (Frequency hopping spread spectrum); deprecated
 - ◆ DSSS (Direct sequence spread spectrum)
 - ◆ Up to 11 Mbps
- 802.11a/g - 2.4 GHz ISM band / 5.0 GHz UNII band
 - ◆ OFDM (Orthogonal frequency domain multiplexing)
 - ◆ Up to 54 Mbps
- 802.11n – 2.4/5.0 GHz bands
 - ◆ Adds MIMO and other tricks to 802.11g
 - ◆ Up to 300-500 Mbps!
- Each backwards compatible with the previous ones



IEEE 802.11b

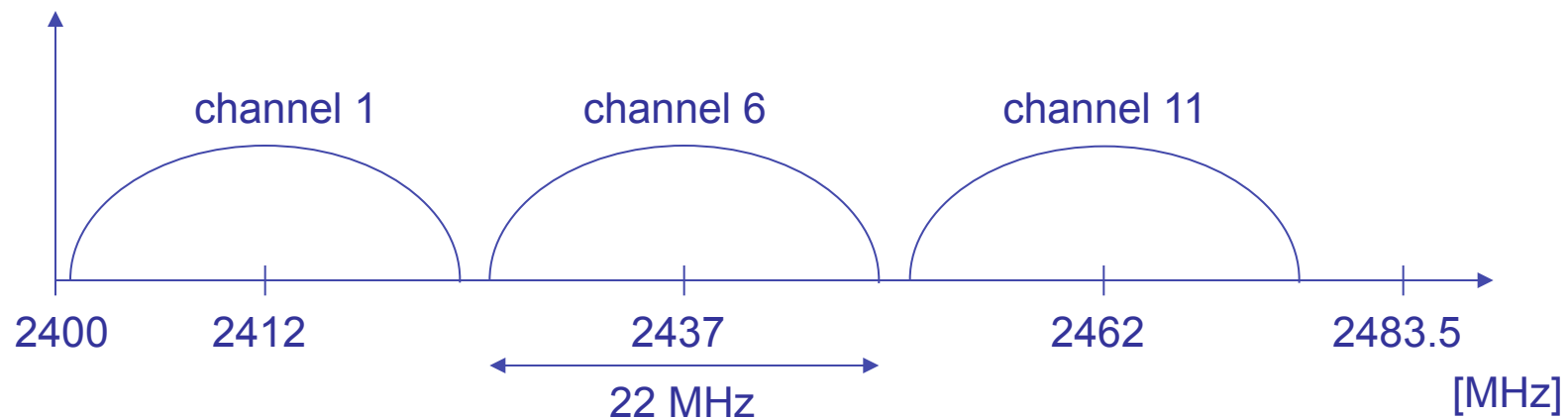
- Data rate
 - ◆ 1, 2, 5.5, 11 Mbit/s
 - ◆ User data rate max. approx. 6 Mbit/s
- Transmission range
 - ◆ 300m outdoor, 30m indoor
 - ◆ Max. data rate ~10m indoor
- Frequency
 - ◆ Free 2.4 GHz ISM-band

802.11b Physical Channels



- 12 channels available for use in the US
 - ◆ Each channel is 20+2 MHz wide
 - ◆ Only 3 orthogonal channels
 - ◆ Using any others causes interference

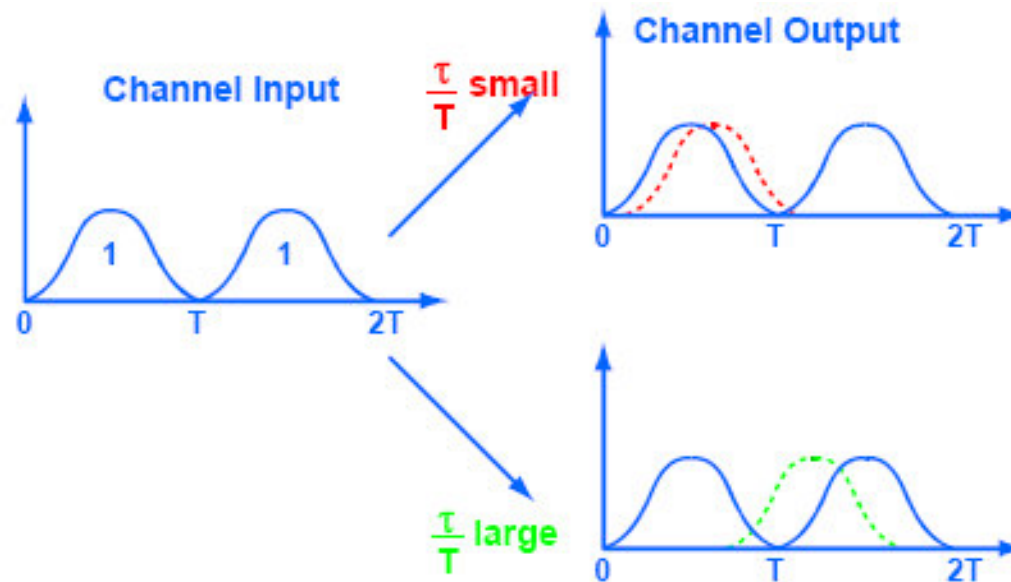
US (FCC)/Canada (IC)





Multipath Interference

- RF signals bounce off of objects (e.g., walls)
 - ◆ Reflected signals travel different distances to receiver
 - ◆ Difference in distance leads to difference in delay

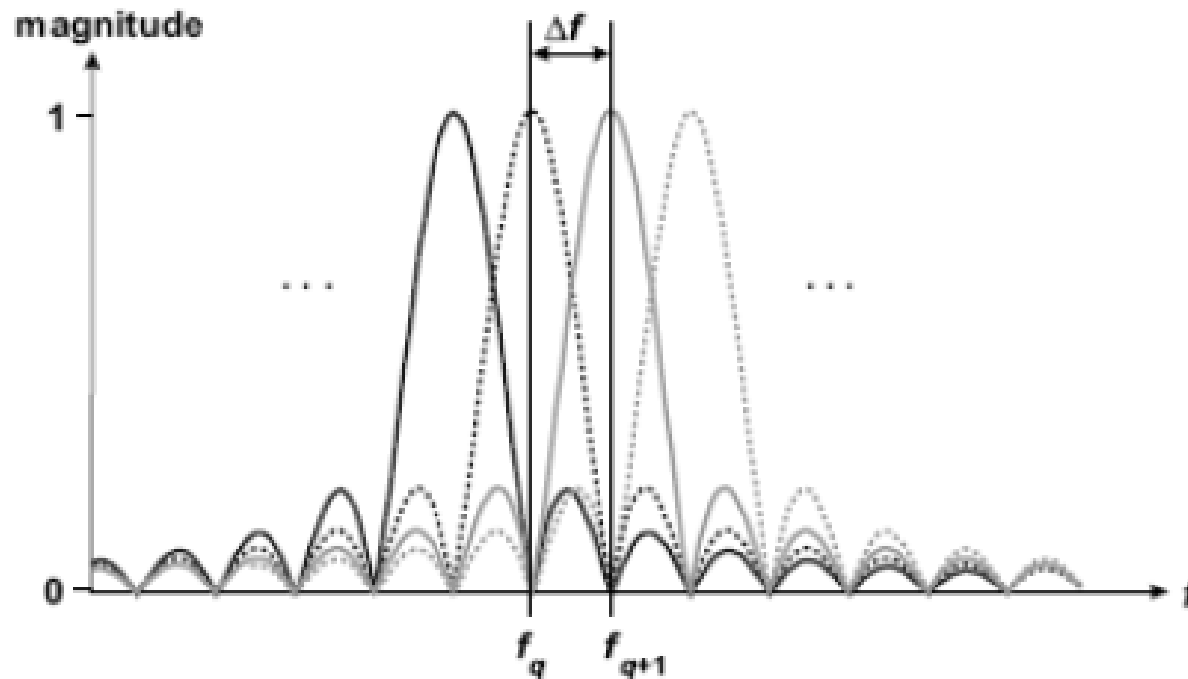


- Limits effective modulation rate in 802.11b



Avoiding ISI: OFDM

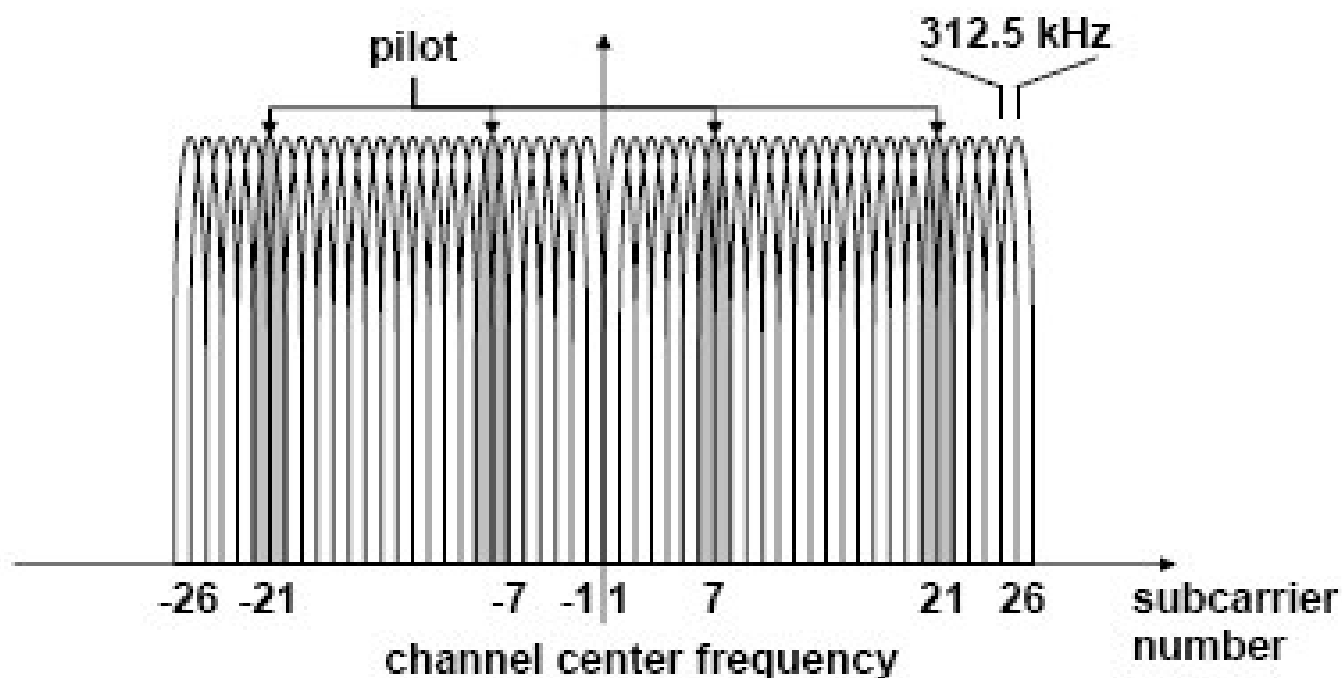
- Break data up into multiple separate streams
 - ◆ Transmit each stream independently on different frequency
 - ◆ Pack frequencies so that they are orthogonal





802.11a/g/n OFDM PHY

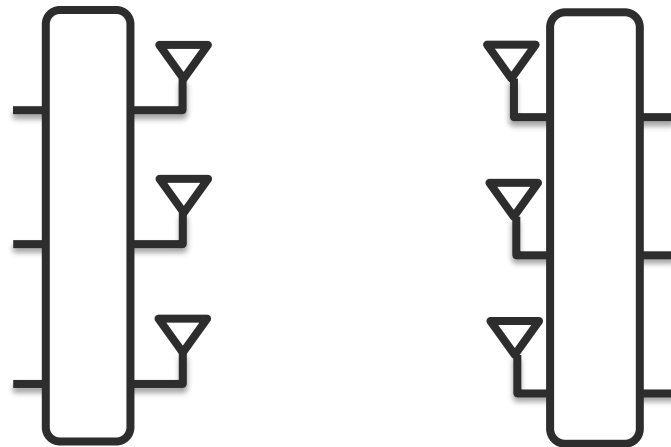
- Each 20-MHz channel divided into 50 subcarriers
 - ♦ Subcarriers spaced appropriately, 4 used as “pilots”



802.11n: MIMO



- Use multiple physical antennae simultaneously
 - ◆ Spatial multiplexing: split data cross antennae
 - ◆ Space-Time Block Coding: same data, encoded differently
 - ◆ Transmit beamforming: steer the signal toward the receiver



Carrier Sense Multiple Access



CSMA: listen before transmit

- If channel sensed idle: transmit entire packet
- If channel sensed busy, defer transmission
 - ◆ Persistent CSMA: retry immediately with probability p when channel becomes idle (may cause instability)
 - ◆ Non-persistent CSMA: retry after random interval
- But what about collisions?

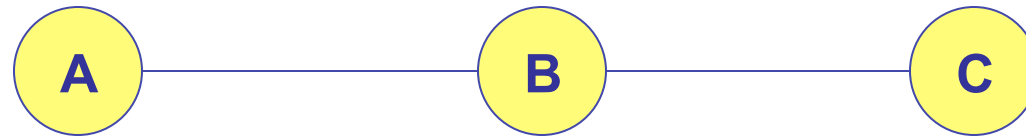


CSMA/CA

- Impossible to hear collision w/half-duplex radio
- Wireless MAC protocols often use **collision avoidance** techniques, in conjunction with a **(physical or virtual) carrier sense** mechanism
- Collision avoidance
 - ◆ Nodes negotiate to reserve the channel.
 - ◆ Once channel becomes idle, the node waits for a randomly chosen duration before attempting to transmit.

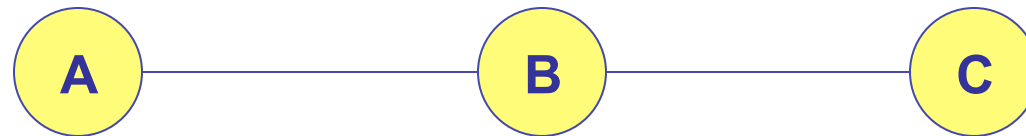


Hidden Terminal Problem



- B can communicate with both A and C
- A and C cannot hear each other
- Problem
 - ◆ When A transmits to B, C cannot detect the transmission using the **carrier sense** mechanism
 - ◆ If C transmits, collision will occur at node B
- Solution
 - ◆ Hidden sender C needs to defer

RTS/CTS (MACA)



- When A wants to send a packet to B, A first sends a **Request-to-Send (RTS)** to B
- On receiving RTS, B responds by sending **Clear-to-Send (CTS)**, provided that A is able to receive the packet
- When C overhears a CTS, it keeps quiet for the duration of the transfer
 - ♦ Transfer duration is included in both RTS and CTS



Backoff Interval

- **Problem:** With many contending nodes, RTS packets will frequently collide
- **Solution:** When transmitting a packet, choose a backoff interval in the range $[0, CW]$
 - ◆ CW is contention window
- Wait the length of the interval when medium is idle
 - ◆ Count-down is suspended if medium becomes busy
 - ◆ Transmit when backoff interval reaches 0
- Need to adjust CW as contention varies



MILD Algorithm in MACAW

- MACAW uses exponential increase linear decrease to update CW
 - ◆ When a node successfully completes a transfer, reduces **CW** by 1
 - ◆ In 802.11 CW is restored to CW_{min}
 - ◆ In 802.11, CW reduces much faster than it increases
- MACAW can avoid wild oscillations of CW when many nodes contend for the channel



Cute Hack

- We can use CTS to reserve the channel for ourselves
 - ◆ Don't use RTS/CTS handshake, just back half
 - ◆ Called a CTS-to-self, simply transmit CTS before our packet
- Doesn't solve hidden terminal, but does squelch
 - ◆ Means stations don't need to be able to decode data frame
- 802.11g uses CTS-to-self to operate w/802.11b
 - ◆ 11g stations always send a CTS before sending packets encoded in a way (OFDM) that 11b stations can't decode
- Much more efficient than full RTS/CTS

Challenge: Reliability

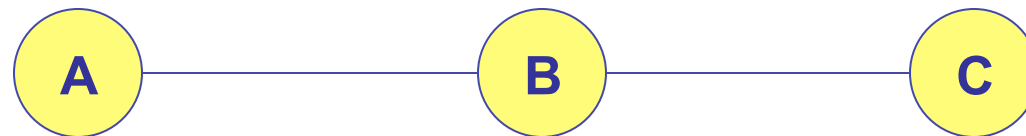


- Wireless links are prone to errors. High packet loss rate detrimental to transport-layer performance.
- Mechanisms needed to reduce packet loss rate experienced by upper layers



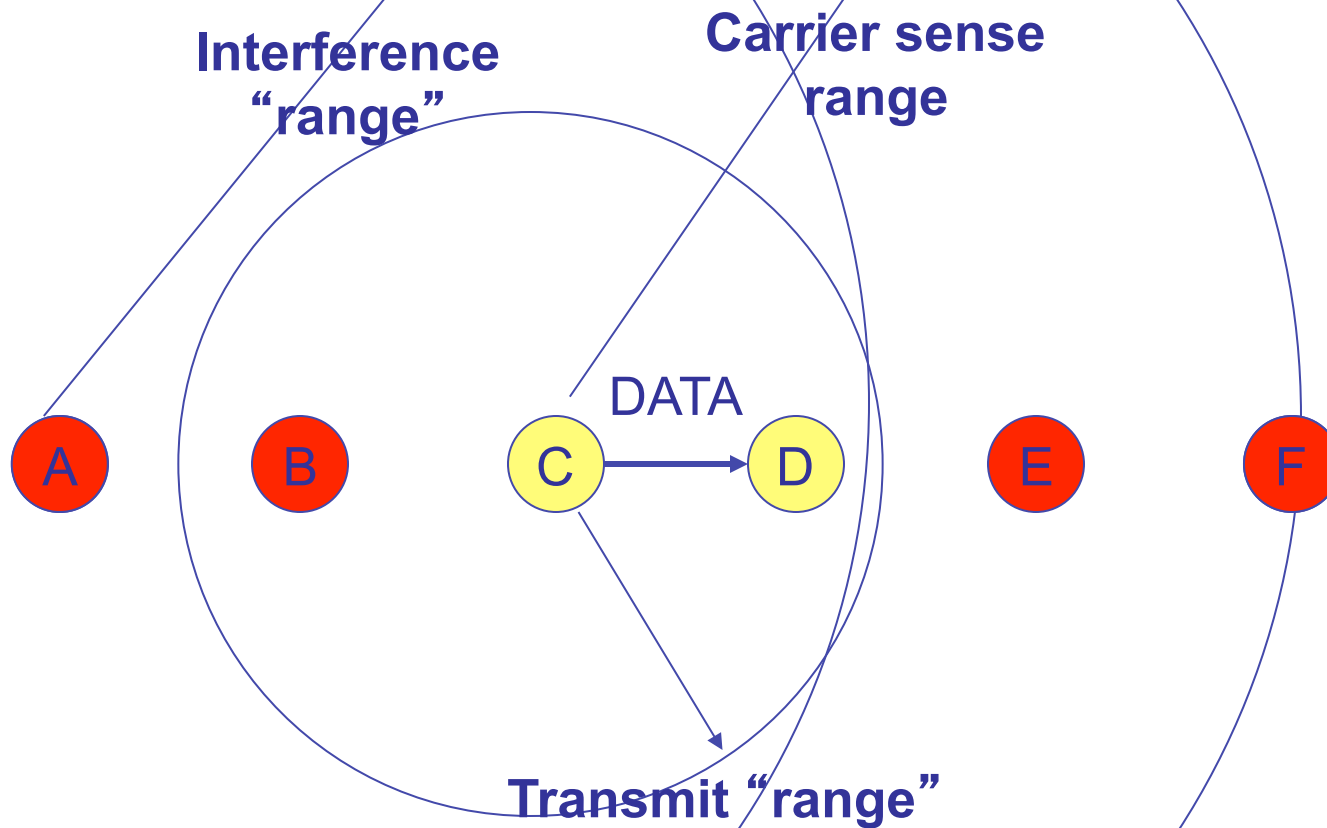
Link-layer ARQ

- When B receives a data packet from A, B sends an Acknowledgement (ACK) to A.
- If node A fails to receive an ACK, it will retransmit the packet





Non-symmetric ranges





Other MACAW Features

- Fairness: Normally, each node wins the channel with equal probability
 - ◆ Nodes with multiple streams should be more aggressive
 - ◆ Abandoned in 802.11. Why?
- Conservative collision avoidance
 - ◆ Use a Data Sending (DS) packet to reserve the channel
 - ◆ 802.11 uses different length intervals and the NAV
- Request-for-Request-to-Send
 - ◆ Assume carrier sense range far larger than transmission range



802.11 MAC Modes

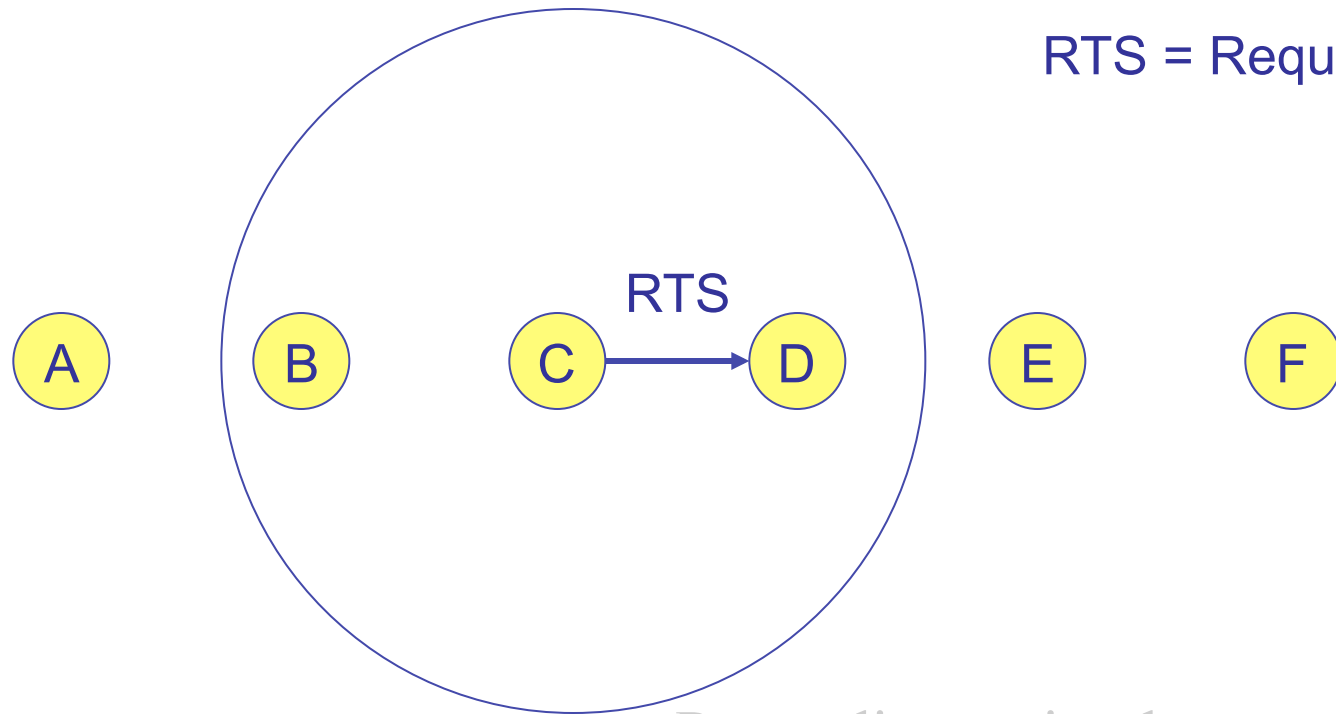
- Distributed Coordination Function (DCF) CSMA/CA
 - ◆ collision avoidance via randomized “back-off” mechanism
 - ◆ minimum distance between consecutive packets
 - ◆ ACK packet for acknowledgements (not for broadcasts)
- DCF w/ RTS/CTS
 - ◆ Distributed Foundation Wireless MAC
 - ◆ avoids hidden terminal problem
- Point Control Function (PCF) - *optional*
 - ◆ Access point polls terminals according to a list
 - ◆ We’ re not going to discuss...

IEEE 802.11 DCF



- DCF is **CSMA/CA** protocol
 - ◆ Uses a Network Allocation Vector (NAV) to implement collision avoidance
- DCF suitable for multi-hop ad hoc networking
- Optionally uses RTS-CTS exchange to avoid hidden terminal problem
 - ◆ Any node overhearing a CTS cannot transmit for the duration of the transfer
- Uses ACK to provide reliability

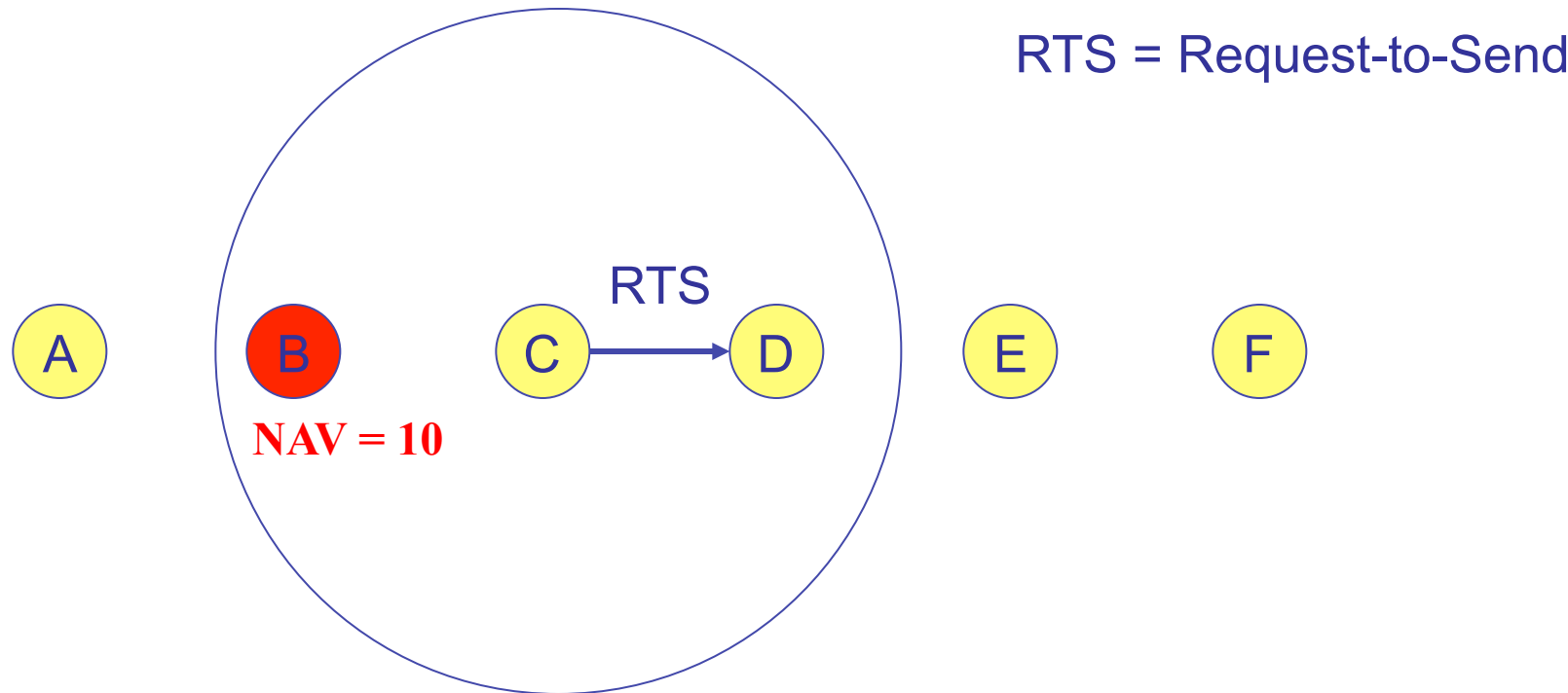
IEEE 802.11



RTS = Request-to-Send

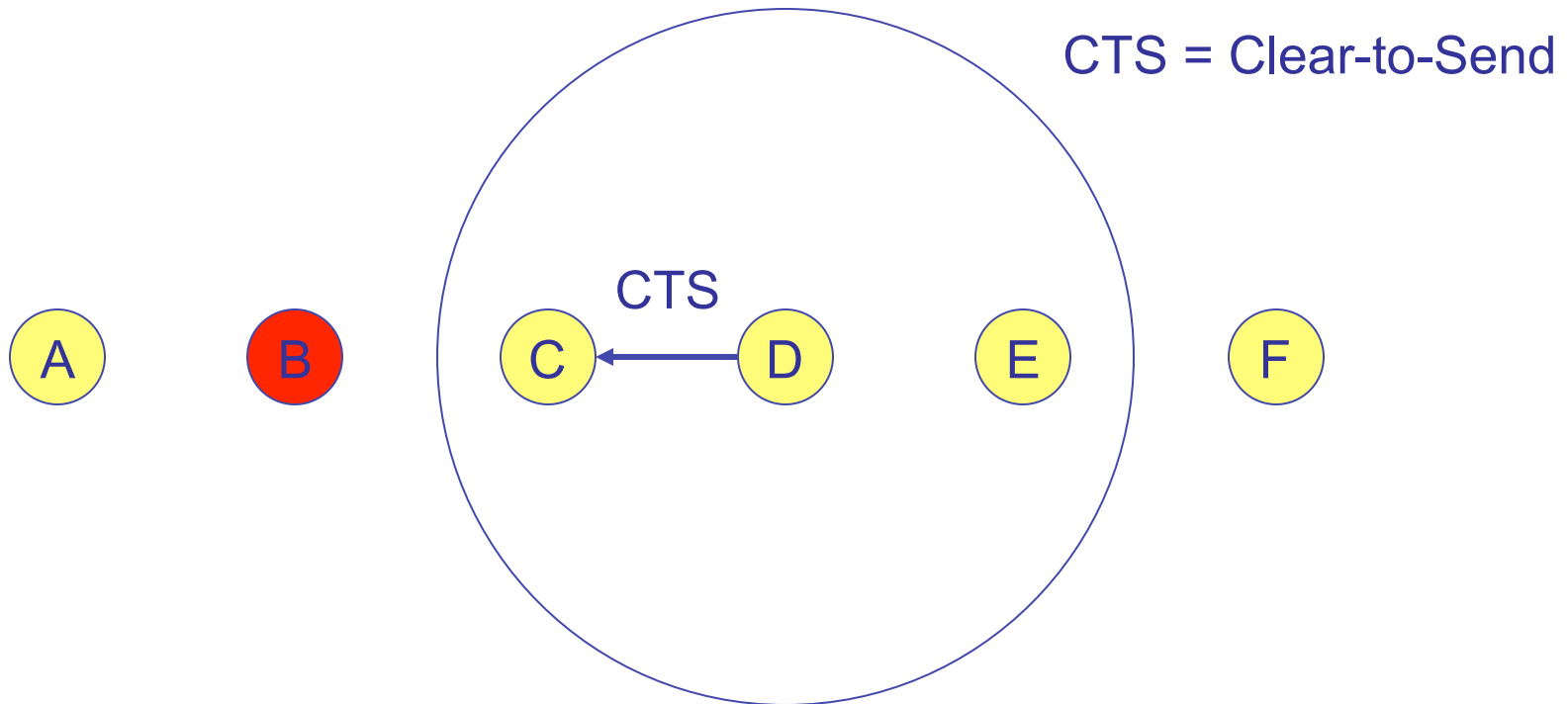
Pretending a circular range

IEEE 802.11

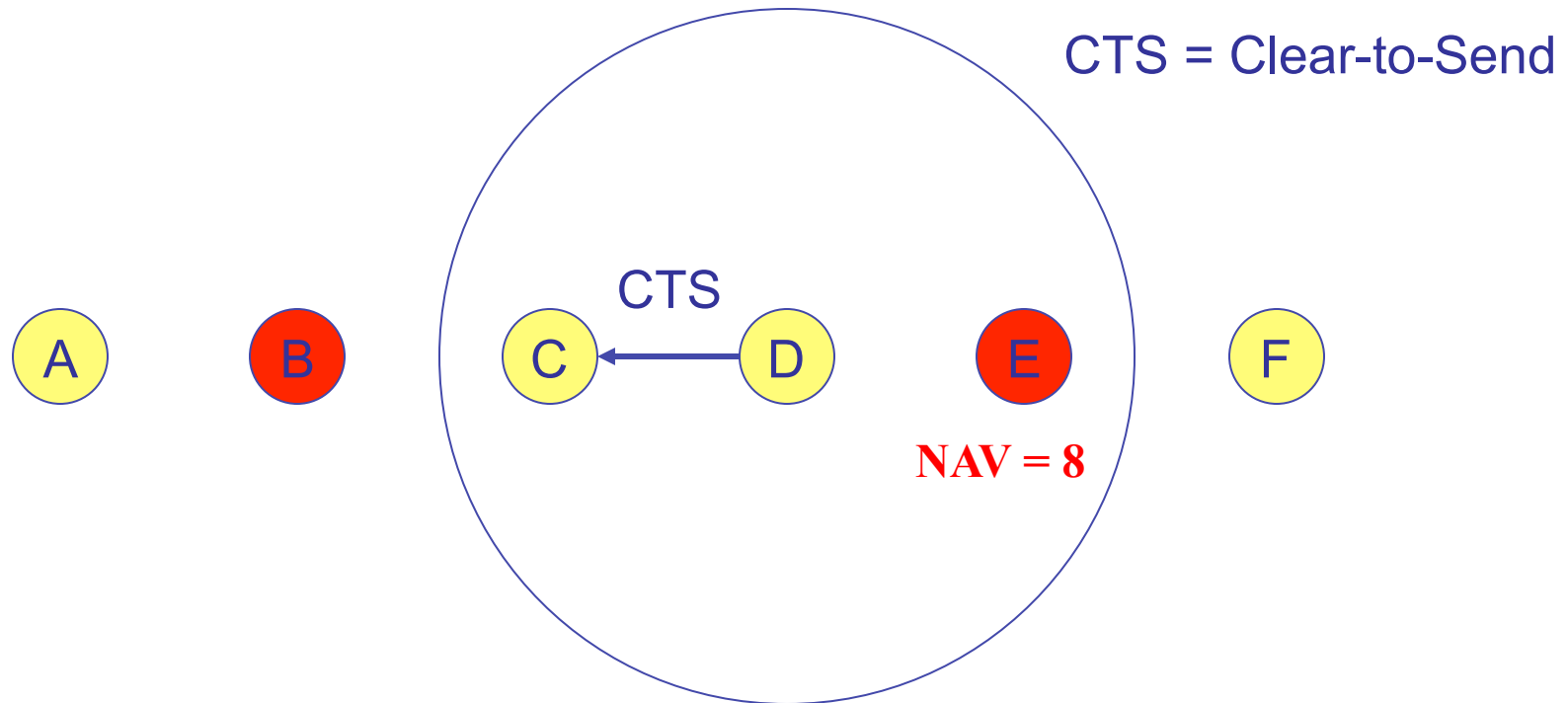


NAV = remaining duration to keep quiet 26

IEEE 802.11



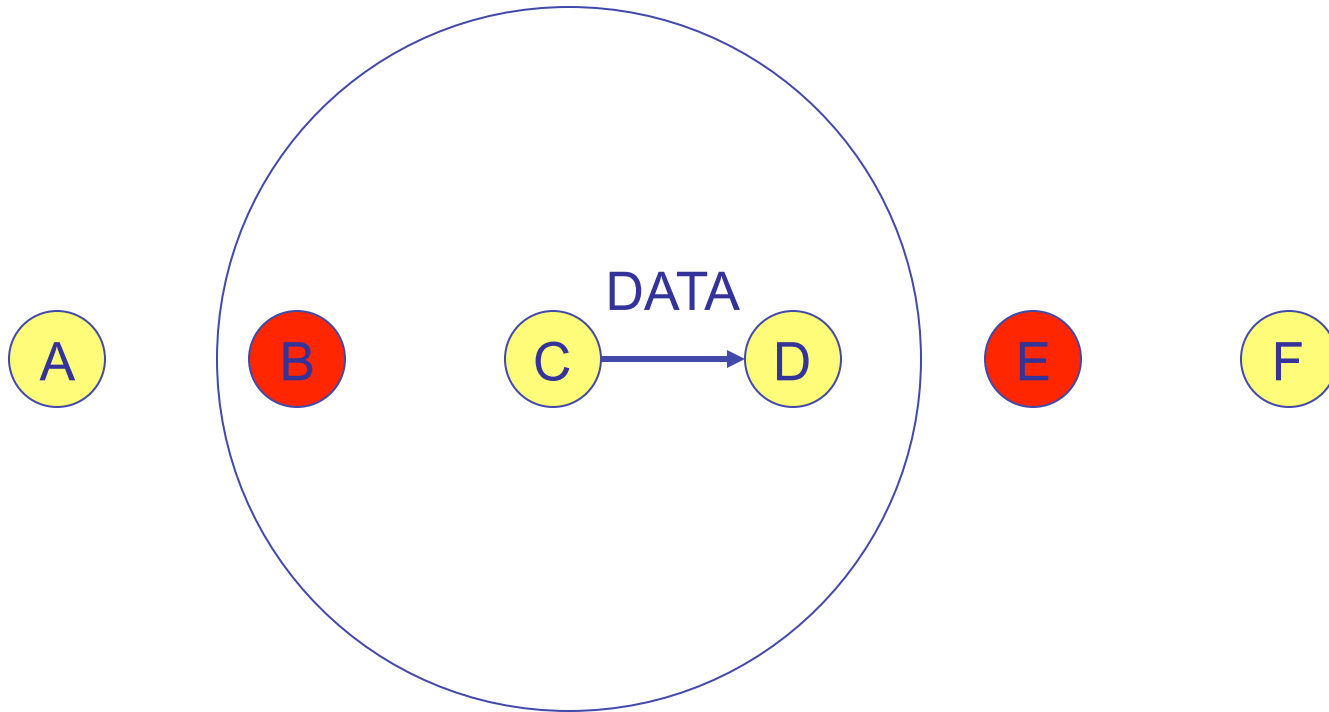
IEEE 802.11



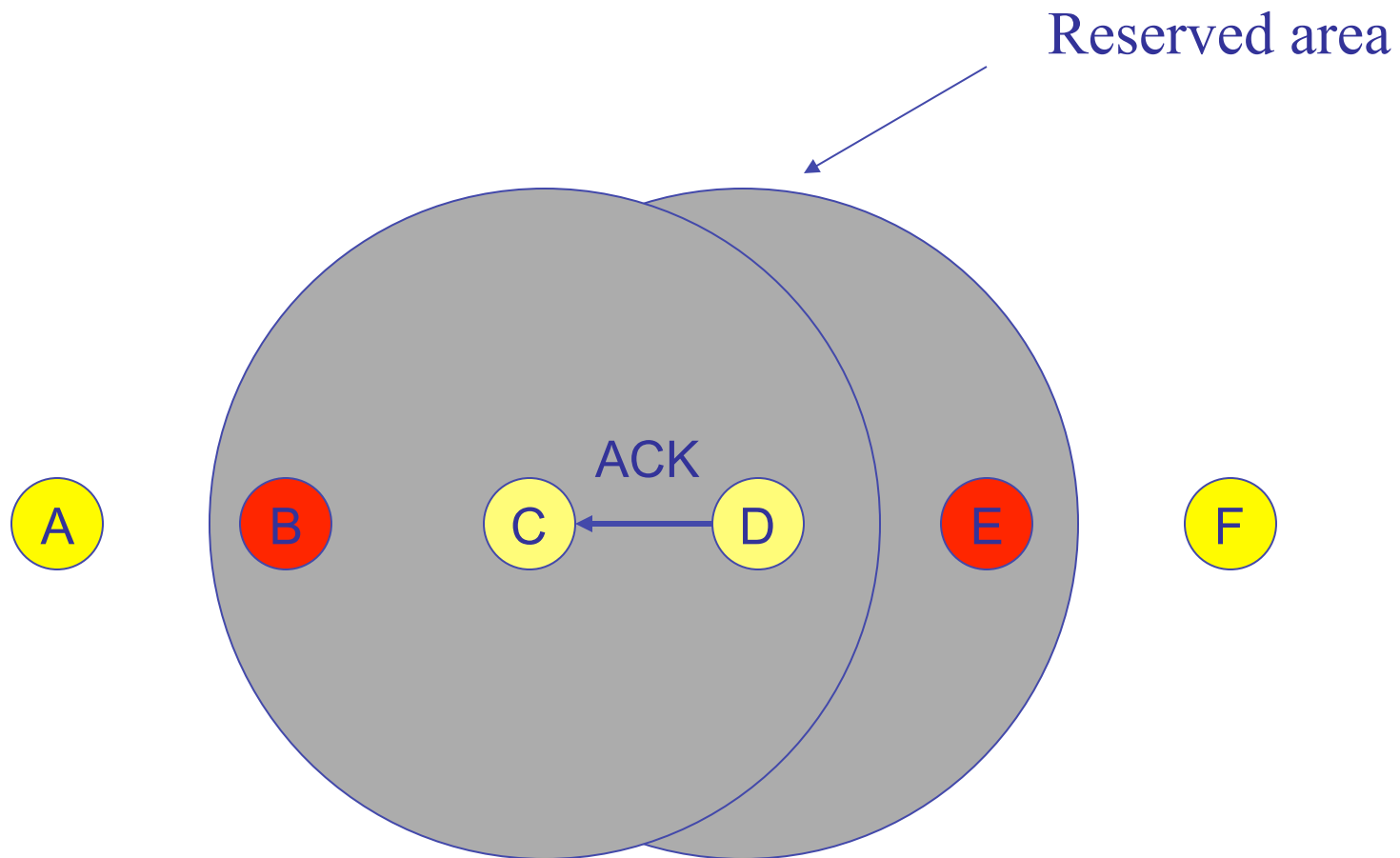
IEEE 802.11



- **DATA** packet follows CTS. Successful data reception acknowledged using **ACK**.



IEEE 802.11





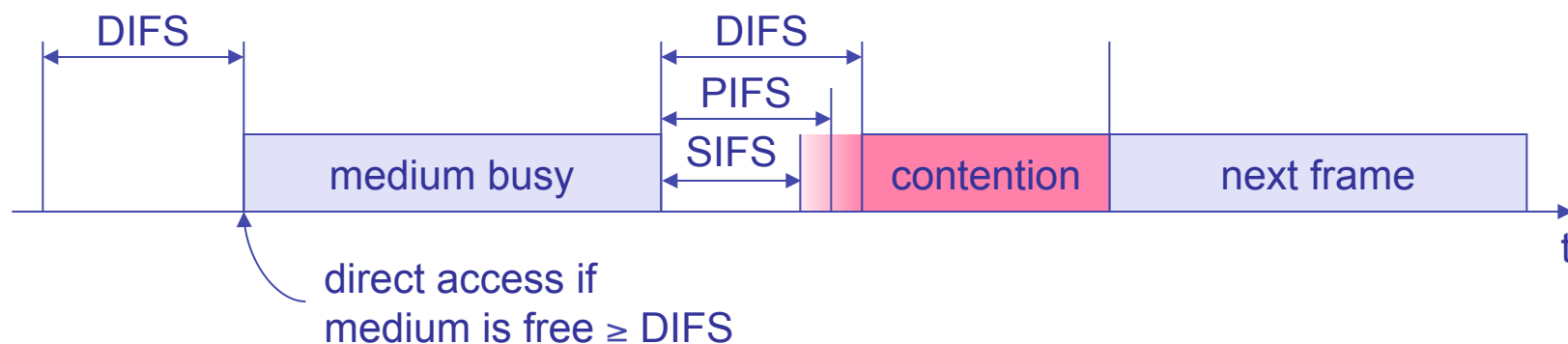
Binary Exponential Backoff in DCF

- When a node fails to receive CTS in response to its RTS, it increases the contention window
 - ♦ CW is doubled (up to an upper bound)
 - ♦ More collisions \rightarrow longer waiting time to reduce collision
- When a node successfully completes a data transfer, it restores CW to CW_{min}

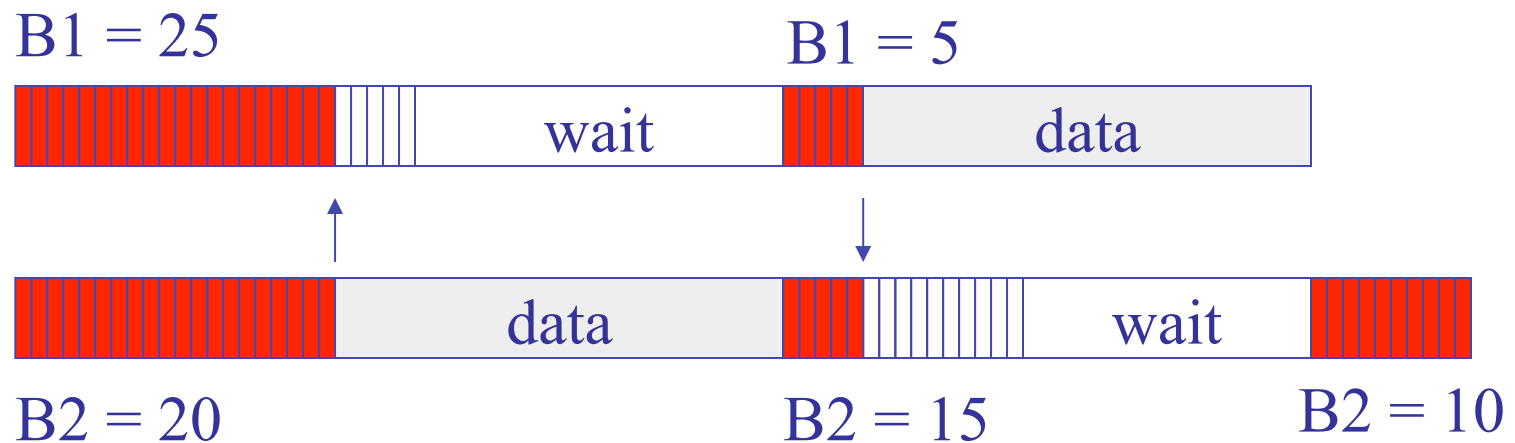


802.11 Backoffs

- SIFS (Short Inter Frame Spacing)
 - ♦ highest priority, for ACK, CTS, polling response
- PIFS (PCF IFS)
 - ♦ medium priority, for time-bounded service using PCF
- DIFS (DCF, Distributed Coordination Function IFS)
 - ♦ lowest priority, for asynchronous data service



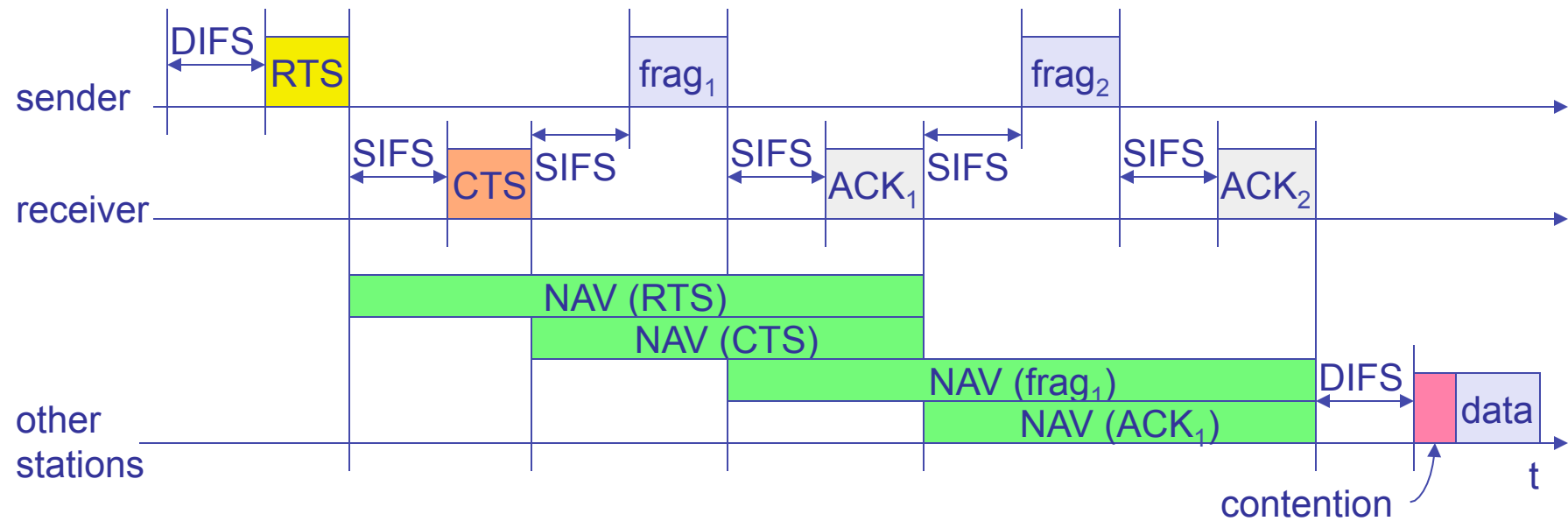
DCF Example



$cw = 31$

**B1 and B2 are backoff intervals
at nodes 1 and 2**

Fragmentation



802.11 - MAC management



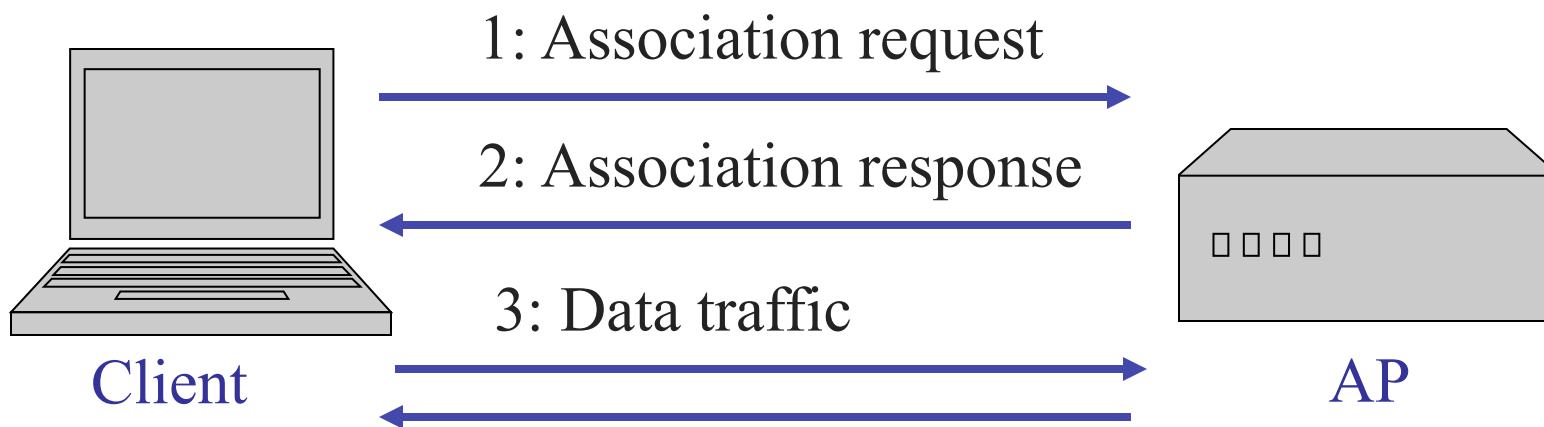
- Association/Reassociation
 - ◆ integration into a LAN
 - ◆ roaming, i.e. change networks by changing access points
 - ◆ scanning, i.e. active search for a network
- Power management
 - ◆ sleep-mode without missing a message
 - ◆ periodic sleep, frame buffering, traffic measurements



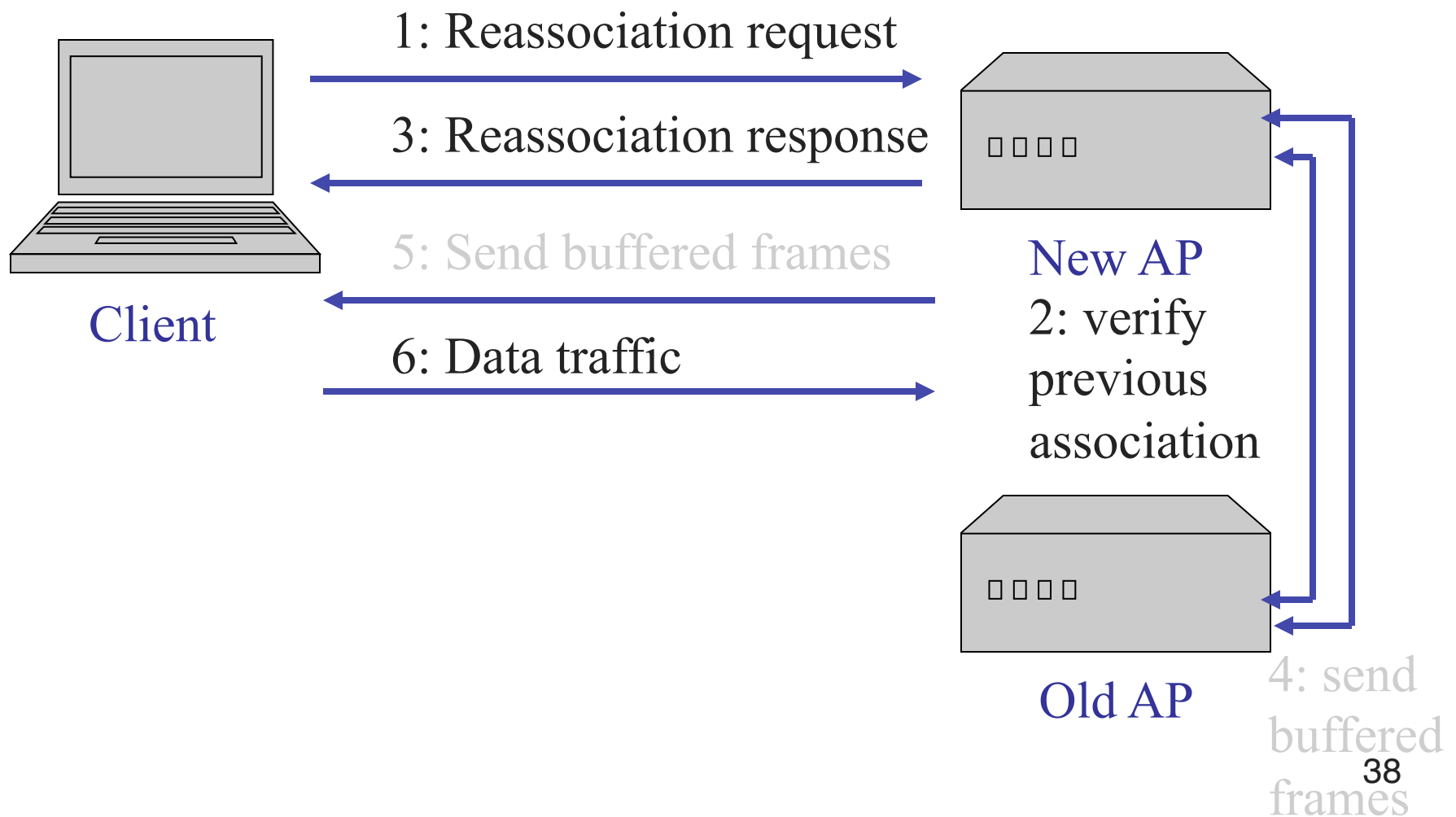
Scanning

- Goal: Find a network to connect
- Passive scanning
 - ◆ Not require transmission
 - ◆ Move to each channel, and listen for Beacon frames
- Active scanning
 - ◆ Require transmission
 - ◆ Move to each channel, and send Probe Request frames to solicit Probe Responses from a network

Association in 802.11



Reassociation in 802.11





802.11 - Roaming

- No or bad connection? Then perform:
- Scanning
 - ♦ scan the environment, i.e., listen into the medium for beacon signals or send probes into the medium and wait for an answer
- Reassociation Request
 - ♦ station sends a request to one or several AP(s)
- Reassociation Response
 - ♦ success: AP has answered, station can now participate
 - ♦ failure: continue scanning
- AP accepts Reassociation Request
 - ♦ signal the new station to the distribution system
 - ♦ the distribution system updates its data base (i.e., location information)
 - ♦ typically, the distribution system now informs the old AP so it can release resources



Power management

- Idea: switch the transceiver off if not needed
- States of a station: sleep and awake
- Timing Synchronization Function (TSF)
 - ◆ stations wake up at the same time
- Infrastructure
 - ◆ Traffic Indication Map (TIM)
 - » list of unicast receivers transmitted by AP
 - ◆ Delivery Traffic Indication Map (DTIM)
 - » list of broadcast/multicast receivers transmitted by AP

802.11 PSM

